Structurally synthesized binary decision diagrams
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Binary graph is an oriented acyclic connected graph with root and two leaves (sinks) – 0 and 1. Every intermediate node $v$ has two successors: $high(v)$ and $low(v)$.

Binary graphs are a skeleton of binary decision diagrams (BDD): BDD is a binary graph in which intermediate nodes are labelled by propositional variables, the sink 1 is labelled by the truth value 1 and the sink 0 by the truth value 0. We denote the label of the node $v$ by $label(v)$. 
Let $D$ be a decision diagram with variables $x_1, \ldots, x_n$. Every vector $\alpha \in B^n$ activates a path $p(\alpha) = p_0, \ldots, p_k$ in $D$ from the root to a terminal node: if $\alpha \vdash \text{label}(v_i)$, then $v_{i+1} = \text{high}(v_i)$ else $v_{i+1} = \text{low}(v)$. We call the path an 1-path (0-path), if its last element is the sink 1 (0). The Boolean function $f_D(x_1, \ldots, x_n)$, represented by $D$ is defined: $f(\alpha) = 1$ iff the path, activated by $\alpha$, is a 1-path.
Let $G$ and $E$ be two binary graphs and $v$ an intermediate node in $G$.

**Definition.** A *superposition* of $E$ into $G$ instead of $v$ ($G_{v \leftarrow E}$) is a graph, which we receive by deleting $v$ from $G$ and redirecting all edges, pointing to $v$, to the root of $E$, all edges of $E$ pointing to sink 1 to the node $\text{high}(v)$, and all edges pointing to the sink 0 to the node $\text{low}(v)$.
We define three elementary binary graphs: A, C and D.

a) $A$

\begin{tikzpicture}
  
  \node (u) at (0,0) [shape=circle,draw] {$u$};
  \node (1) at (1,0) [shape=square,draw] {$1$};
  \node (0) at (1,-1) [shape=square,draw] {$0$};

  \draw[->] (u) -- (1);
  \draw[->] (1) -- (0);

\end{tikzpicture}
b) $C$

Diagram:

- A circle labeled $u$ connected to a circle labeled $v$.
- The circle $v$ is connected to a rectangle labeled $1$.
- A rectangle labeled $0$ is connected to the circle $v$. 
c) $D$
Definition. We define inductively a family of superpositional graphs:

1° An elementary graph $A$, $C$ or $D$ is a superpositional graph.

2° If $G$ and $E$ are superpositional graphs and $v$ is a node in $G$, then a superposition of $E$ instead of $v$ into $G$, $G_{v\leftarrow E}$, is an superpositional graph.
Definition. *high-path (low-path)* is the path $p = (v_0, \ldots, v_k)$, where $v_{i+1} = \text{high}(v_i)$ ($v_{i+1} = \text{low}(v_i)$) for every $i : 0 \leq i < k$.

**Definition.** Node $v$ in the superpositional graph is a *final* node if $\text{high}(v) = 1$ and $\text{low}(v) = 0$. 
Theorem 1  If $S$ is an superpositional graph, then:

1. $S$ is a planar graph.
2. There exists a high-path from every node to the sink 1.
3. There exists a low-path from every node to the sink 0.
4. $S$ has exactly one final node.
5. There exists a directed path through all intermediate nodes.
6. For every pair of intermediate nodes $u, z$ there exists a directed path from $u$ to $z$ or from $z$ to $u$. 
**Definition.** *Structurally synthesized binary decision diagram* for a formula $F$, $\mathcal{D}(F)$, is a superpositional graph, defined inductively according to the structure of $F$:

1° If $F$ is a literal $l$, then $\mathcal{D}(F)$ is a graph $A$, where the root is labelled by $l$.

2° If $F = P \& R$ then $\mathcal{D}(F)$ is a graph $C_{u \leftarrow \mathcal{D}(P), v \leftarrow \mathcal{D}(R)}$

3° If $F = P \lor R$ then $\mathcal{D}(F)$ is a graph $D_{u \leftarrow \mathcal{D}(P), v \leftarrow \mathcal{D}(R)}$
Theorem 2 Propositional formula $F$ and its SSBDD $D(F)$ are representing the same Boolean function.
Not every path of the $\mathcal{D}(F)$’s superposition graph can be activated. Therefore we define a *consistent* path of the SSBDD $\mathcal{D}(F(x_1, \ldots, x_n))$ as a path which can be activated by some assignment $(\alpha_1, \ldots, \alpha_n) \in B^n$. 
The *satisfiability problem* for SSBDD can be formulated: does there exist a consistent path from the root to the sink 1? It is obvious, that the number of intermediate nodes of the SSBDD $D(F)$ is the number of occurrences of variables in the formula $F$. It follows, that the satisfiability problem is NP-complete and counting the number of true assignments is $\#P$ complete for SSBDD representation of Boolean functions.
It is obvious, that the conjunction of literals in a 1-path in $\mathcal{D}(F)$ is a term of the DNF for $F$. Next theorem (first proved in \cite{?}) says, that we can remove from the terms of the DNF all literals which are false for the corresponding assignment.

**Definition** Positive term of the path $p = (v_1, \ldots, v_l)$ is a conjunction of literals of the nodes $v_i$, followed by the node $\text{high}(v_i)$:

$$\bigwedge_{0 \leq i < l} \text{label}(v_i)$$

$$v_{i+1} = \text{high}(v_i)$$

**Theorem 3** Let $F$ be a propositional formula and $\mathcal{D}(F)$ its SSBDD. Disjunction of positive terms of all 1-paths of $\mathcal{D}(F)$ is a DNF for a propositional formula $F$. 


Logic circuit and its EPF

\[
y = c_y \bar{c}_y = c_y \lor \bar{c}_y = x_{6,c,y} x_{7,3,c,y} \lor d_{e,y} b_{e,y} = \\
= x_{6,c,y} x_{7,3,c,y} \lor (\bar{x}_{1,d,e,y} \lor \bar{a}_{d,e,y})(x_{5,b,e,y} \lor x_{7,2,b,e,y}) = \\
= x_{6,c,y} x_{7,3,c,y} \lor (x_{1,d,e,y} \lor x_{2,a,d,e,y} x_{7,1,a,d,e,y})(x_{5,b,e,y} \lor x_{7,2,b,e,y}) = \\
= x_{6} x_{7,3} \lor (x_{1} \lor x_{2} x_{7,1})(x_{5} \lor \bar{x}_{7,2})
\]
Logic circuit and its SSBDD

\[ x_6 \overline{x}_3 \lor (\overline{x}_1 \lor x_2 x_7) (\overline{x}_5) \lor \overline{x}_7 \]
Fault simulation. Stuck-at fault model

Stuck-at fault model:
1. Only one circuit line is faulty.
2. The faulty line is permanently stuck to 0 or 1.

Stuck-at fault simulation for a circuit line $l$ is equivalent to calculating the Boolean derivative for variable $x_l$. 
Fault simulation. SSBDD model

We say that fault stuck-at $\neg \alpha_i$ in node $v_j$, where $\text{label}(v_j)$ is $x_i$, is covered by a vector $\alpha_1, \ldots, \alpha_i, \ldots, \alpha_n$ iff all of the three conditions are satisfied:

1. There exists an activated path from the root node to $v_j$.
2. There exists an activated path from $\text{high}(v_j)$ to sink $1$.
3. There exists an activated path from $\text{low}(v_j)$ to sink $0$. 
Fault simulation. SSBDD model

\[ \alpha = (1, 0, 0, 1, 0, 0) \]

(i.e. \( x_1 = 1, x_2 = 0, x_3 = 0, x_4 = 1, x_5 = 0, x_6 = 0 \))
Fault simulation. SSBDD model

Number of collapsed and SSBDD faults in ISCAS'85

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Fault simulation. SSBDD model

Correlation between speed up and gates/subcircuits ratio
Conclusions

• A special case of BDDs called Structurally Synthesized BDDs (SSBDD) was proposed.
• SSBDD representations for both, Boolean formulae and digital logic circuits were discussed.
• The paper showed that the special properties of SSBDDs are highly useful in fault modeling and simulation for digital circuits.
• Experiments showed a speed-up of 2 to 7 times in comparison to traditional gate-level descriptions.